

# Cover Attacks on Trace-Zero groups

Jasper Scholten

COSIC,  
Katholieke Universiteit Leuven

Joint work with Claus Diem

# Introduction

Let  $C$  be a smooth projective curve of genus  $g = g(C)$  defined over a finite field  $K = \mathbb{F}_{q^n}$ .

Let  $\text{Cl}^0(C)$  be the degree-0-divisor class group of  $C$ , i.e.:

$$\text{Cl}^0(C) = \frac{\{K\text{-rational divisors of degree } 0\}}{\{\text{principal divisors}\}}$$

Facts:

- $\text{Cl}^0(C) \cong \text{Jac}(C)(K)$
- $\#\text{Cl}^0(C) \approx q^{ng}$

# The Discrete Logarithm Problem

Assume that  $\text{cl}^0(C)$  can not be embedded in the additive group or multiplicative group of a relatively small finite field using pairings. (This is a condition on  $\#\text{cl}^0(C)$ .)

Let  $\ell$  be the largest prime divisor of  $\#\text{cl}^0(C)$ .

Methods for solving the DLP in  $\text{cl}^0(C)$ :

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 $\mathcal{O}\left(g^5 q^{2 - \frac{4}{2g+1}} + \epsilon\right)$ .
- High-genus Index Calculus (Adleman, DeMarrais, Huang; Gaudry, Enge): subexponential with  $\alpha = \frac{1}{2}$ .

# Cover Attacks

(aka Weil descent attacks, or Weil restriction attacks).

Dates back to Frey's talk at ECC 1998.

Recall that  $C$  is defined over  $K = \mathbb{F}_{q^n}$ .

Let  $k = \mathbb{F}_q$ .

Let  $X$  be a curve defined over  $k$ .

Let  $X_K = X \otimes_k K$  be the curve  $X$  considered over the field  $K$ .

# Cover Attacks

A map

$$c : X_K \rightarrow C$$

induces a map

$$\phi : \text{Cl}^0(C) \xrightarrow{c^*} \text{Cl}^0(X_K) \xrightarrow{\text{norm}} \text{Cl}^0(X)$$

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Two conditions:

- The order- $\ell$  subgroup of  $\text{cl}^0(C)$  is not in  $\ker \phi$ .
- $g(X) \approx ng(C)$ .

# Cover Attacks

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- Gaudry, Hess, Smart et al. (2000): Construction in characteristic 2, based on Artin-Schreier theory.
- Some odd-characteristic cases using Artin-Schreier and Kummer theory (Hess, Thériault, Diem).
- This talk: Constructions for trace-zero groups, based on Galois theory of curves.

# Trace-zero groups

Suppose  $C$  can already be defined over the small field  $k$ :

$$C = C_K = C_k \otimes_k K$$

for some  $C_k$  defined over  $k$ .

Let

$$T := \ker (\mathrm{Cl}^0(C_K) \rightarrow \mathrm{Cl}^0(C_k))$$

be the trace-zero group.

$$\#T \approx \frac{(\#K)^g}{(\#k)^g} = q^{(n-1)g}.$$

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Examples

- Koblitz curves ( $n$  is large).
- Lange, Avanzi: efficient arithmetic for  $g(C) = 2$  and  $n = 3$ .

We concentrate on the case that  $n$  is small.

# First Approach

Suppose

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$$\begin{array}{ccc} \mathrm{Cl}^0(C_K) & \xrightarrow{c^*} & \mathrm{Cl}^0(X_K) \\ & \searrow \phi & \downarrow \text{norm} \\ & & \mathrm{Cl}^0(X_k) \end{array}$$

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So

$$T \subset \ker \phi.$$

Not what we want...

# Twists

Modify this construction as follows:

Suppose

$$\tau : X_k \longrightarrow X_k$$

is an automorphism of order  $n$ .

One can twist  $X_k$  by  $\tau$  over  $K/k$ . Call this twist  $X_k^\tau$ . This is a curve defined over  $k$ , and there is a canonical isomorphism between the base-changed curves  $\psi : X_K^\tau \rightarrow X_K$ .

The set of  $k$ -rational points  $X_k^\tau(k)$  can be identified with

$$\{P \in X_k(K) \mid \tau^{-1}(P) = \text{Frob}_q(P)\}.$$

Consider the modified cover

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We have the induced map  $\phi' : \text{Cl}^0(C_K) \longrightarrow \text{Cl}^0(X_k^\tau)$ .

The corresponding Trace-zero subgroup  $T$  should not be mapped to 0. A necessary condition for this is that

$$c \circ \tau \neq c$$

# How to construct such $X$ ?

With Galois theory!

Assume  $n$  is prime. Let  $k(C_k)$  be the function field of  $C_k$ .

Let  $f \in k(C_k)$  be a function of degree  $n$ .

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Consider the extension of fields

$$\begin{array}{c} F \quad , \\ | \\ k(C_k) \\ | \quad n \\ k(f) \end{array}$$

where  $F$  is the Galois closure. So  $F/k(f)$  is a Galois extension whose Galois group  $G$  contains an element  $\tau$  of order  $n$ .



# Curves instead of fields

Assume that  $k$  is the exact constant field in  $F$ .  $F$  is then the function field of some curve  $X$  defined over  $k$  with an automorphism  $\tau$  of order  $n$ .

$$\begin{array}{c} X \\ \downarrow \\ C_k \\ \downarrow n \\ \mathbb{P}_k^1 \end{array}$$

$g(C)$  and  $g(X)$  are determined in terms of the ramification indices by the Hurwitz formula:

$$2g(C) - 2 = -2n + \sum_{P \in C_k(\bar{k})} (e_P - 1).$$

# Galois theory of curves

Assume for now that we work over  $\mathbb{C}$ . Let  $P_1, \dots, P_r \in \mathbb{P}^1(\mathbb{C})$  denote  $r$  distinct points.

Let  $\pi = \pi_1(\mathbb{P}^1(\mathbb{C}) \setminus \{P_1, \dots, P_r\})$  denote the fundamental group. One knows that

$$\pi \cong \langle s_1, \dots, s_r \mid s_1 \cdot s_2 \cdot \dots \cdot s_r = 1 \rangle.$$

There is a one to one correspondence between the following sets:

$$\{\text{covers } f : C \rightarrow \mathbb{P}^1 \text{ unramified outside } \{P_1, \dots, P_r\}\} / \sim$$

and

$$\{\text{finite sets with transitive } \pi\text{-action}\} / \sim$$

# Galois theory of curves

Let  $S = \{1, \dots, n\}$  be a finite set with transitive  $\pi$ -action.

Denote the action of  $s_i$  by  $\sigma_i$ .

The corresponding cover  $f : C \rightarrow \mathbb{P}^1$  satisfies:

- $\deg(f) = n$ .
- The ramification indices above  $P_i$  are exactly the cycle lengths of  $\sigma_i$ .

Hence the genus  $g(C)$  is determined by the group-theoretic data.

# Galois theory of curves

Let

$$X \longrightarrow C \longrightarrow \mathbb{P}^1$$

be the Galois closure. Then

- The Galois group  $G$  of  $X \rightarrow \mathbb{P}^1$  is  $\langle \sigma_1, \dots, \sigma_r \rangle \subset S_n$ .
- The unique ramification index above  $P_i$  is  $\text{ord}(\sigma_i)$ .

So the genus  $g(X)$  is also determined by this data.

# Remarks on this construction

The Galois theory of curves over  $\overline{\mathbb{F}}_q$  is the same as over  $\mathbb{C}$  as long as the characteristic does not divide the group order  $\#G$ .

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Theorem If  $G$  is the group of affine linear transformations  $x \mapsto ax + b$  of  $\mathbb{Z}/n\mathbb{Z}$ , then  $\ker \phi'$  is annihilated by  $n$ .

So for big  $\ell$ , the order- $\ell$  subgroup of  $\text{Cl}^0(C)$  is not mapped to zero.

# Summary of examples

Examples for this theorem.

- $n = 3, G = S_3, g(C) = 2, g(X) = 6$ : Can be done for every curve  $C$ .

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- $n = 3, G = S_3, g(C) = 2, g(X) = 4$ : We have an explicit family of such curves  $C$ .
- $n = 7, G$  is the group of affine linear transformations of  $\mathbb{Z}/7\mathbb{Z}$ .  $g(C) = 1, g(X) = 8$ : We have explicit 1-parameter families.

# Examples

Example for  $g(C) = 2$ ,  $n = 3$ ,  $g(X) = 6$ .

Take  $\sigma_1 = (1\ 2\ 3)$ ,  $\sigma_2 = (3\ 2)$ ,  $\sigma_3 = \dots = \sigma_7 = (1\ 2)$ .

Given  $C_k$ , one can construct the function  $f$  as follows:

Take a point  $P \in C_k(k)$ , and take  $f$  from the Riemann-Roch space  $\mathcal{L}(3P) \setminus \mathcal{L}(2P)$ . Then  $f : C_k \rightarrow \mathbb{P}_k^1$  has ramification index 3 above  $\infty$ , and in general there will be 6 other ramification points, all with index 2.

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Security comparison:

Pollard- $\rho$ :  $(\log q)^2 q^2$

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Index calculus is faster for large key sizes

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Example for  $g(C) = 2$ ,  $n = 3$ ,  $g(X) = 5$ .

Take  $\sigma_1 = \sigma_4 = (1\ 2\ 3)$ ,  $\sigma_2 = \sigma_5 = (3\ 2)$  and  $\sigma_3 = \sigma_6 = (2\ 1)$ .

Explicit construction of  $f$ :

Let  $\iota : C \rightarrow C$  denote the hyperelliptic involution.

Suppose  $\text{cl}^0(C_k)$  has a 3-torsion divisor class that is represented by a divisor class of the form  $P_1 + P_2 - O - \iota(O)$  then  $3P_1 - 3\iota(P_2)$  is the divisor of a function  $f$ . Generically, this  $f$  has the required ramification.

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One can construct an explicit family such that the curve  $X$  is hyperelliptic.



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Example for  $g(C) = 2$ ,  $n = 3$ ,  $g(X) = 4$ .

Take  $\sigma_1 = (1\ 2\ 3\ 4\ 5)$ ,  $\sigma_2 = (1\ 2\ 3\ 5\ 4)$  and  $\sigma_3 = (1\ 3)(2\ 4)$ .

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Explicit families can be constructed using Kummer theory.

Example for  $g(C) = 1$ ,  $n = 7$ ,  $g(X) = 8$ .

Take  $\sigma_1 = (1\ 6)(2\ 5)(3\ 4)$ ,  $\sigma_2 = (1\ 7)(2\ 6)(3\ 5)$ ,

$\sigma_3 = (1\ 2\ 4)(3\ 6\ 5)$  and  $\sigma_4 = (2\ 5\ 3)(4\ 6\ 7)$ .

Explicit families can be constructed using Kummer theory.

# Are there other examples?

- $n = 5$ ,  $G = A_5$ ,  $g(C) = 1$ ,  $g(X) = 4$ : There is one such curve:  $y^2 = x^3 + 3165x - 31070$ .

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- There are other isolated examples in Diem's thesis.
- A complete classification is work in progress, joint work with Avanzi.